**Plan Bouquet based Techniques for Variable Sized Databases**

Achint Chaudhary

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**Abstract**

OLAP applications require a certain set of queries to be fired on database with the change in certain constants in queries. For optimal execution of these queries, query optimizer does select an optimal strategy known as the query execution plan. These choices are based on cardinality estimates of various predicates that differ from actual cardinality values encountered during execution. Due to this reason, optimizer given plans led to high inflation in actual execution cost as compared to predicted during optimization.  
  
An altogether different approach for query processing is proposed in 2014, named Plan Bouquet. Basis of which is selectivity discovery at run-time by repeated execution of multiple plans. This technique provides strong bounds independent of data distribution.

Plan Bouquet on other hand is not robust against large updates in the database. This work focuses on providing incremental algorithms that can use past information about plan bouquet and iso-cost contours, to provide further robust execution without incurring entire compilation overhead of plan bouquet.

**Sec.1 Introduction**

Optimal Plan from Optimizer comes from various structural choices of logical and physical operators for query execution. These choices are based on the cost of each operator which is calculated using number of tuples it will process known as . Cardinality normalized in range of [0, 1] is known as throughout our analysis.

These selectivity values are estimated before query execution based on some statistical models used in classical cost-based optimizers. An entirely different approach based on run-time selectivity discovery is proposed called Plan Bouquet, which provides for first time theoretical strong bounds on worst-case performance as compared to optimal performance possible from all the available plan choices.

For each given query, predicates having the potential of selectivity error contributes as a dimension in . The set of Optimal plans over the entire range of selectivity values in is called . POSP is generated by asking optimizer's optimal plans at various selectivity values using Selectivity injection module.

A subset of POSP is identified as Plan bouquet, which is obtained by the intersection of plans trajectory with

Iso-cost surfaces, each of which is placed at some ratio proportion () of cost from the previous surface.

Since these executions form geometric progress, the total cost of which is can be also derived using the sum of geometric progression. The figure below shows the performance of bouquet w.r.t to optimal oracle performance.



Fig 1. Cost incurred by Oracular vs Bouquet Execution

In above figure, various plans up to actual selectivity value are executed. Each plan has a limit provided by the next iso-cost surface. This yields total execution cost of

This leads to Sub-Optimality (ratio of incurred cost to optimal cost) of bouquet approach as

This value is minimized using , which provides theoretical worst case bound of 4 times the optimal execution time.

**Sec. 2 Notations, Notes & Assumptions**

**2.a Notations used in this work**

|  |  |
| --- | --- |
| **Notation** | **Description** |
|  | Query Predicate |
|  | Set of all Query Predicates |
|  | Set of Actually Known Predicates |
|  | Set of Error Prone Predicates |
|  | Set of Trivial predicates |
|  | EPP Selectivity Space |
|  | Dimension of ESS |
|  | Cost of Optimal Plan at selectivity |
|  | Scaling Factor of Predicate in Database |
|  | Minimum Selectivity on Predicate |
|  | Resolution on Each axis of ESS |
|  | Number of Iso-cost Contours |
|  | Cost Ratio of Iso-cost contours |
|  | Discretized interval for each axis of ESS |
|  | Set of Optimal Plans over entire ESS |
|  | Cost of plan P at selectivity location q in ESS |
|  | Cardinality of predicate p at location q with database size scale s |
|  | Difference in Selectivity values |
|  | Ratio of Selectivity values |
| β | Worst case slope of Plan Cost Function |
| α | Tolerance of contour thickening |

**2.b Notes**

**2.b.1 Interplay of Selectivity & Cardinality**

Selectivity is the fraction of tuples out of maximum possible tuples that can come out of a query predicate. Notation of selectivity is devised to make study of ESS independent of cardinality values.

**2.b.2 Distribution of Selectivity values on axis of ESS**

is and will be discretized at finite resolution during initial compilation with number of points . Each axis denoting selectivity for each of predicate in will have points. One choice that will matter is whether to take these points in Arithmetic progression (AP) or Geometric progression (GP) in interval .

GP gives more focus on lower selectivities (near origin) since the plans and cost rapidly changes in that region of ESS. Choice of GP also helps in designing efficient incremental bouquet algorithm.

is minimum considered selectivity, sooner we will see bounds on this also under some assumptions.

**2.b.3 Trivial Predicates**

These are predicates never seen when dealing with a fixed size database, as they always have maximum selectivity of 1.0. Example of this is scanning of entire relation without any filter.

These predicates are useful when comparing old and new database with reference to each other.

**2.c Assumptions**

**2.c.1 Plan Cost Monotonicity (PCM)**

This assumption implies that if location dominates location in selectivity component of each predicate, processing more tuples will have more cost

**2.c.2 Perfect Cost Model of Optimizer**

This assumption states that poor choices of plan come only from cardinality estimation error of optimizer and not from cost model itself. While, we have assumed perfect cost model of optimizer, an optimizer with bounded cost model will also work well. Improving, which is an orthogonal problem, one work on static tuning [4], proves most actual cost values to lie in 30% of estimated cost values after tuning.

**2.c.3 Axis Parallel Concavity (APC)**

This assumption is on Plan Cost Function () which is not just monotonic but exhibit a weak form of in their cost trajectories. For 1D , is said to be concave if for any two selectivities locations from and any below condition holds

Generalizing to D dimensions, a PCF is said to be if the function is concave along every axis-parallel 1D segment of .

Which simply states that each PCF should be concave along every vertical and horizontal line in the ESS

Further, an important and easily provable implication of the exhibiting APC is that the corresponding , which is the infimum of the PCFs, also satisfies APC. Finally, for ease of presentation, we will generically use concavity to mean APC in the remainder of this work.

**2.c.4 are only query predicates**

For present analysis, we have considered that all of query predicates are Error-prone, there is also no trivial predicate, which means each relation has some filter applied over it.

Rational behind this assumption is that, if we supply same selectivity value to both old and new database, their outputs will be of different cardinalities. Also, we need to determine change of scale for and , in that case, maybe re-compilation of bouquet is needed.

**Sec. 3 Challenges in database size change**

Plan bouquet is suitable for canned queries as compilation overhead of is amortized over repeated invocations. Under situation of change of database size, placements of ideal contours can be totally different from existing contours built for old database size, which may result that old bouquet plans are totally different from new bouquet contours.

Under above mentioned conditions it seems that a re-compilation will be needed. While, in this work we will focus on how efforts done in earlier compilation can be utilized for incremental compilation, and also to perform performance guarantees similar to Plan Bouquet Approach.

**Sec. 4 Present Directions**

**4.a Using Old Contours**

Compilation phase of Bouquet based approach needs , this much cost cannot be afforded if database size changes frequently, also what we care is change of scale of each predicate as compared to what it was when bouquet was initially compiled.

Placement of contours in bouquet works like we first create a contour of and keep placing other contours from cost ratio less than the cost of previously created contour.

In this section we prove that we can:

1. Use previous complete contours
2. Extend previous incomplete contour into extended
3. Create new contours in extended region of when needed

Focusing on a 1-Dimensional version, where a predicate is scaled up, scale change can be evaluated as

With assumption, we can view predicate on new database, as scaled up version of predicate on old database by extending maximum selectivity up to .

From now onwards, our entire formulation takes old database as reference, because bouquet is compiled on that.

Formula for number of contours

Using above expression, number of contours in old and new database will be

If is extended to , in any dimensional ESS, the last contour is always a point, a single optimal plan at that location just needs an optimizer call, so that last contour need not be saved during incremental compilation.

IF [ is an INTEGER ] :

ELSE:

A screenshot of a cell phone

Description automatically generated

Fig 2. Placement of Ideal re-drawn contours vs Extended contours of initial bouquet

**4.b Proof of Performance Guarantee**

Having evaluated , which is number of new contours to place at cost ratio from last contour of previously compiled bouquet ( in Fig 2). These extended contours are presented as in Fig 2, also plan under entire re-compilation are shown as

Also, if last contour of has cost value lesser than ( which is ), then an additional last contour is constructed which will always contain an optimal plan at maximum selectivity location.

If this contour exist, it will have cost ratio which is less than from last constructed contour out of contours built during incremental compilation.

Proof will be carried out in two cases.

**Case 1.** When selectivity discovery of using bouquet execution can be completed up to contours.

Above case satisfies same Suboptimality condition as original bouquet and will give performance guarantee of 4 using .

**Part 2.** When selectivity discovery of using bouquet execution will need execution of

Since, we have shown in figure that last contour if it exists will have lesser cost ratio from previous compiled iso-cost contour as . Under same consideration expression for suboptimality is

Sine has upper bound of , when is substituted in above expression and upper bound of is used for . Suboptimality in that case is also upper bounded by 4.

Hence, we can state that incremental bouquet compilation under assumption of , will provide same worst-case performance guarantee.

This shows that in the region of for which Bouquet is already compiled no redundant effort is required.

**4.c Incremental Bouquet Compilation Approach**

**4.c.1 Constructing New contours**

Given which is number of extra contours that needs to be drawn in new . Portion of which is newly constructed should be searched for these contours and bound on those efforts are shown in later sub-section.

**4.c.2 Completing existing contours**

It can be case that existing contours, needs to be completed in extended . Contours that have either intersection on boundary beyond anti-diagonal should be completed during incremental compilation.

In both constructing of new contours and completing new one w.r.t new, we will require an effort of

With above formulation, we should try to minimize that volume of new should not be too high w.r.t old .

**4.d Number of Extra Points Needed on each axis**

Ratio for Geometric progression of selectivity values of old data base and new database to be created are

Extending same relation to new scale of predicate and total points

Re-arranging to find , we get

To solve this equation further, we need to find what is, this minimum selectivity possible should also have some lower bound otherwise there are too large jumps on selectivity values on axis in values close to maximum selectivity of for incremental bouquet.

To derive lower bound on ℇ, we can consider worst case of concavity assumption, which is linear

Using extreme lower bound value of , we get

Extra points needed on axis will be , which will be

This bound on Number of points needed is independent of RES. While if we have done through Arithmetic progression-based range of selectivities on each axis, it will be

**4.d Efforts needed for incremental compilation**

Since done with previous section, which claims that linear number of more points will be needed with even exponential increase in database size, incremental compilation will need

Which is much less than, what we get using arithmetic progression for selectivity values on each axis for arbitrarily large scale-up on each predicate axis of .

**5 Experiments**

**6 Conclusion and Future Work**

Each predicate axis in will demand extra points that are only logarithm of scale change of that predicate from old to new database, which is also independent of taken to discretize .

Given proof of concept is not able to handle the change of scale in , for now, to handle this we can go with a conservative assumption of , and consider also as comes with curse of dimensionality. we can apply dimensionality reduction techniques [3], if not anything else to handle .

Also, there can be some trivial predicates () within a query, like a simple scan of a relation without any filter will result different cardinalities on different sized databases, and hence will impact optimal plans generated for on both old and new sized database. To take care of this, Trivial predicates should be created and combined with in future experiments.

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